# Landscapes and Other Art Forms. 

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## Introduction to Elementary Landscapes

What is a landscape?

- Many different intuitive definitions
- A mathematical formalism of the search space of a combinatorial optimization problem

Definition: a landscape is a tuple $(X, N, f)$

$$
\begin{array}{ll}
\text { A set of states } & X \\
\text { A neighborhood operator } & N: X \mapsto \mathcal{P}(X) \\
\text { A fitness function } & f: X \mapsto \mathbb{R}
\end{array}
$$

## Introduction: a landscape


$1,4,5,6,3,2$
$X$ set of states

## Introduction: a landscape



$$
\begin{aligned}
X & \text { set of states } \\
N: X \mapsto \mathcal{P}(X) & \text { neighborhood operator }
\end{aligned}
$$

## Introduction: a landscape



$$
\begin{aligned}
X & \text { set of states } \\
N: X \mapsto \mathcal{P}(X) & \text { neighborhood operator } \\
f: X \mapsto \mathbb{R} & \text { objective function }
\end{aligned}
$$

## Preliminaries

$$
G(X, E)
$$

is the underlying graph induced by $N$.
We assume $G$ is regular with vertices of degree $d$.

$$
\mathbf{A} \in \mathbb{R}^{|X| \times|X|}
$$

is the adjacency matrix of $G$.
If $x_{1}$ and $x_{2}$ are neighbors, $A\left(x_{1}, x_{2}\right)=1$.

$$
\boldsymbol{\Delta}=\mathbf{A}-d \mathbf{I}
$$

is the Laplacian of $G$.

## The Wave Equation: definition 1

On an arbitrary landscape

- $f$ and $N$ are unrelated

On an elementary landscape
The wave equation

$$
\Delta f=\lambda f
$$

- where $\lambda$ is a scalar
- In other words, $f$ is an eigenvector of the Laplacian


## The Wave Equation: definition 1

Average change

$$
\begin{gathered}
\boldsymbol{\Delta} f=(\mathbf{A}-d \mathbf{I}) f=k(\bar{f}-f) \\
\boldsymbol{\Delta} f(x)=\sum_{y \in N(x)}(f(y)-f(x))=k(\bar{f}-f(x))
\end{gathered}
$$

Average value

$$
\begin{aligned}
\underset{y \in N(x)}{\operatorname{avg}}\{f(y)\} & =\frac{1}{d} \sum_{y \in N(x)} f(y) \\
& =f(x)+\frac{1}{d}\left(\sum_{y \in N(x)} f(y)-f(x)\right) \\
& =f(x)+\frac{1}{d} \Delta f(x) \\
& =f(x)+\frac{k}{d}(\bar{f}-f(x))
\end{aligned}
$$

## The Wave Equation: definition 2

$$
f(x)=\sum \text { a subset of "components" }
$$

Starting from average...

$$
\underset{y \in N(x)}{\operatorname{avg}}\{f(y)\}=f(x)+\underset{y \in N(x)}{\operatorname{avg}}\{\text { components in - components out }\}
$$

## Example: TSP under 2-opt



- Components: set of edge weights $w_{i, j}$
- $f(x)=$ sum of edge weights induced by tour $x$
- There are $n(n-1) / 2-n$ weights not in tour $x$
- Average value of components out: $\frac{2}{n} f(x)$
- Average value of components in: $\frac{2}{n(n-3) / 2}\left(\sum w-f(x)\right)$


## The Components and $\bar{f}$

Let $C$ denote the set of components
$0<p_{3}<1$ is the proportion of the components in $C$ that contribute to the cost function for any randomly chosen solution

$$
\bar{f}=p_{3} \sum_{c \in C} c
$$

For the TSP:

$$
\begin{gathered}
\bar{f}=\frac{n}{n(n-1) / 2} \sum_{w_{i, j} \in C} w_{i, j} \\
\bar{f}=\frac{2}{n-1} \sum_{w_{i, j} \in C} w_{i, j}
\end{gathered}
$$

## The Components and $\bar{f}$

$C$ is the set of components (e.g. from a cost matrix)
$x$ is a solution (e.g. a subset of the cost matrix)
Let $(C-x)$ denote the set of components, excluding those in $x$
For a move, we then define:
$0<p_{1}<1$ is the proportion of components in $x$ that change
$0<p_{2}<1$ is the proportion of components in $(C-x)$ that change

## The Components and $\bar{f}$

## Theorem

If $p_{1}, p_{2}$ and $p_{3}$ can be defined for any regular landscape such that the evaluation function can be decomposed into components where $p_{1}=\alpha / d$ and $p_{2}=\beta / d$ and

$$
\bar{f}=p_{3} \sum_{c \in C} c=\frac{\beta}{\alpha+\beta} \sum_{c \in C} c
$$

then the landscape is elementary.

## The Wave Equation: definition 2

$$
\begin{aligned}
\underset{y \in N(x)}{\operatorname{avg}}\{f(y)\} & =f(x)+\frac{2}{n(n-3) / 2}\left(\sum w-f(x)\right)-\frac{2}{n} f(x) \\
& =f(x)+\frac{2}{n(n-3) / 2}((n-1) / 2 \bar{f}-f(x))-\frac{2}{n} f(x) \\
& =f(x)+\frac{(n-1)}{n(n-3) / 2}(\bar{f}-f(x)) \\
& =f(x)+\frac{k}{d}(\bar{f}-f(x))
\end{aligned}
$$

For a 5 city TSP

|  | $a b$ | $b c$ | $c d$ | de | ae | ac | ad | $b d$ | be | ce |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ABCDE | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| ABEDC | 1 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| ABCED | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 |
| ABDCE | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |
| ACBDE | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 1 | 0 | 0 |
| ADCBE | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 0 |

Looking at the neighbors in aggregate.

| ab | bc | cd | de | ae | ac | ad | bd | be | ce |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 |
| 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 |

## Properties



## A General Model for all bounded Pseudo-Boolean Problems

This obviously applies to MAXSAT and NK Landscapes


## Superpositions of Elementary Landscapes

$$
f(x)=f 1(x)+f 2(x)+f 3(x)+f 4(x)
$$

$$
f 1(x)=f 1_{a}(x)+f 1_{b}(x)+f 1_{c}(x)
$$

$$
f 2(x)=f 2_{a}(x)+f 2_{b}(x)+f 2_{c}(x)
$$

$$
f 3(x)=f 3_{a}(x)+f 3_{b}(x)+f 3_{c}(x)
$$

$$
f 4(x)=f 4_{a}(x)+f 4_{b}(x)+f 4_{c}(x)
$$

$$
\varphi^{(1)}(x)=f 1_{a}(x)+f 2_{a}(x)+f 3_{a}(x)+f 4_{a}(x)
$$

$$
\varphi^{(2)}(x)=f 1_{b}(x)+f 2_{b}(x)+f 3_{b}(x)+f 4_{b}(x)
$$

$$
\varphi^{(3)}(x)=f 1_{c}(x)+f 2_{c}(x)+f 3_{c}(x)+f 4_{c}(x)
$$

$$
f(x)=\varphi^{(1)}(x)+\varphi^{(2)}(x)+\varphi^{(3)}(x)
$$

## MAX-3SAT decomposition

MAX-3SAT is a superposition of 3 elementary landscapes
Walsh span of order $p$

$$
\varphi^{(p)}=\sum_{\{i: b c(i)=p\}} w_{i} \psi_{i}
$$

The $p^{\text {th }}$ Walsh span is an elementary landscape

$$
\Delta \varphi^{(p)}=-2 p \varphi^{(p)}
$$

With Thanks to Andrew Sutton!

## MAX-3SAT decomposition

Recall that we can express $f$ as:

$$
f(x)=\sum_{i=1}^{m} \sum_{j=1}^{2^{k}} w_{m(i, j)} \psi_{m(i, j)}(x)
$$

Grouping the Walsh decomposition results in

$$
f(x)=\sum_{p=0}^{3} \varphi^{(p)}(x)
$$

Thus MAX-3SAT is a superposition of 3-elementary landscapes

## This allows us to compute Statistical Summaries

## Properties

Average neighborhood

$$
\begin{aligned}
\sum_{y \in N(x)} f(y) & =f(x)+\frac{1}{d} \sum_{y \in N(x)}(f(y)-f(x)) \\
& =f+\frac{1}{d} \Delta f \\
& =f+\frac{1}{d} \Delta\left(\sum_{p=0}^{3} \varphi^{(p)}\right) \\
& =f-\frac{1}{d}\left(\sum_{p=0}^{3} 2 p \varphi^{(p)}\right)
\end{aligned}
$$

For 3-SAT

$$
f(x)-\frac{1}{d}\left(2 \varphi^{(1)}(x)+4 \varphi^{(2)}(x)+6 \varphi^{(3)}(x)\right)
$$

## This allows us to compute Statistical Summaries

We can compute statistics efficiently over Hyperplanes

We can compute statistics efficiently over Hamming Balls

## Random Walks and Autocorrelation

Remark. We have the following identities:

$$
\langle f, f\rangle=\sum_{i} w_{i}^{2} \quad\left\langle f, \mathbf{T}^{s} f\right\rangle=\sum_{i} \lambda_{i}^{s} w_{i}^{2} \quad\langle\mathbf{1}, f\rangle=w_{0}
$$

Random walk process estimates the following equation

$$
r(s)=\frac{\left\langle f, \mathbf{T}^{s} f\right\rangle-\langle\mathbf{1}, f\rangle^{2}}{\langle f, f\rangle-\langle\mathbf{1}, f\rangle^{2}}
$$

Substitutions...

$$
r(s)=\frac{\sum_{i} \lambda_{i}^{s} w_{i}^{2}-w_{0}^{2}}{\sum_{j} w_{j}^{2}-w_{0}^{2}}=\frac{\sum_{i \neq 0} \lambda_{i}^{s} w_{i}^{2}}{\sum_{j \neq 0} w_{j}^{2}}
$$

## Random Walks and Autocorrelation

This gives exact autocorrelation function

$$
r(s)=\frac{\sum_{i \neq 0} \lambda_{i}^{s} w_{i}^{2}}{\sum_{j \neq 0} w_{j}^{2}}
$$

where $\lambda_{i}=\left(1-\frac{2\langle i, i\rangle}{n}\right)$.

Recall for $k$-SAT all nonzero $w_{i}$ can be computed in $O(m)$ time.

## THANK YOU

Take Home Message:
PROBLEM STRUCTURE MATTERS.

Black Box Optimizers can never match the performance of an algorithm that efficiently exploits problem structure.

But we need only a small amount of information: Gray Box Optimization.

