

Squeeziness

A metric for avoiding fault masking in software testing

- Joint work with

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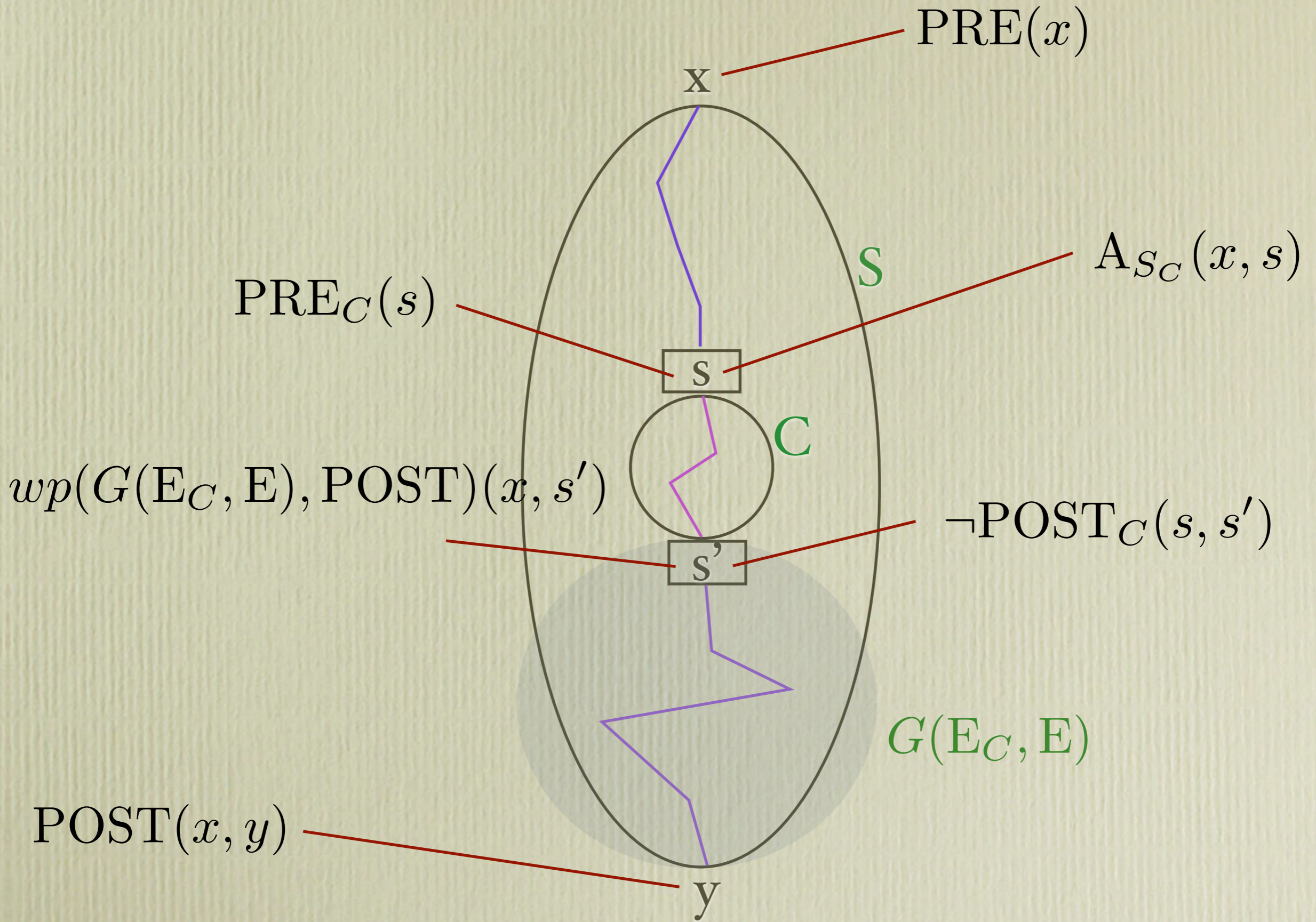
- with help from **Sebastian Hunt** and *Laurie Tratt*

software fault masking

- also called **error masking**
- reduces test set effectiveness
- Error masking condition:

$$\begin{aligned} \exists x, s, s', y . & \text{PRE}(x) \wedge A_{S_C}(x, s) \wedge \text{PRE}_C(s) \\ & \wedge \neg \text{POST}_C(s, s') \wedge wp(\text{G}(\text{E}_C, \text{E}), \text{POST})(x, s') \\ & \wedge \text{POST}(x, y) \end{aligned}$$

Laski et al. '95



example

Intended

`x=x+2;`
`if (x>0)`
 `x=x%4;`
 `else x=x;`

“oracle”

output

`t1:x==1`
`t2:x==-3`

Unintended

`x=3*x;`
`if (x>0)`
 `x=x%4;`
 `else x=x;`

input

`t1:x==3`
`t2:x==-5`

output

`t1:x==1`
`t2:x==-15`

collisions and state abstraction

- `(x > 0) == true; x % 4: collisions`
- **also:** oracle may examine only part of the state
- execution path plus oracle identify good and bad states

Domain to Range Ratio

- collisions *necessary*, not *sufficient*, for fault masking
- [Woodward and al-Khanjari (2000)] observed fault masking associated with domain to range ratio
- “loss of information measure” $|D|/|R|$

information theoretic view

Treat the input space and the output space for a program as random variables: I and O

Oracle's Observation
of Output

Information in a random variable

$$\mathcal{H}(X) = - \sum_{x \in X} p(x) \log_2 p(x)$$

Loss of information from running program \mathbf{P}

$$\mathcal{H}(I) - \mathcal{H}(O)$$

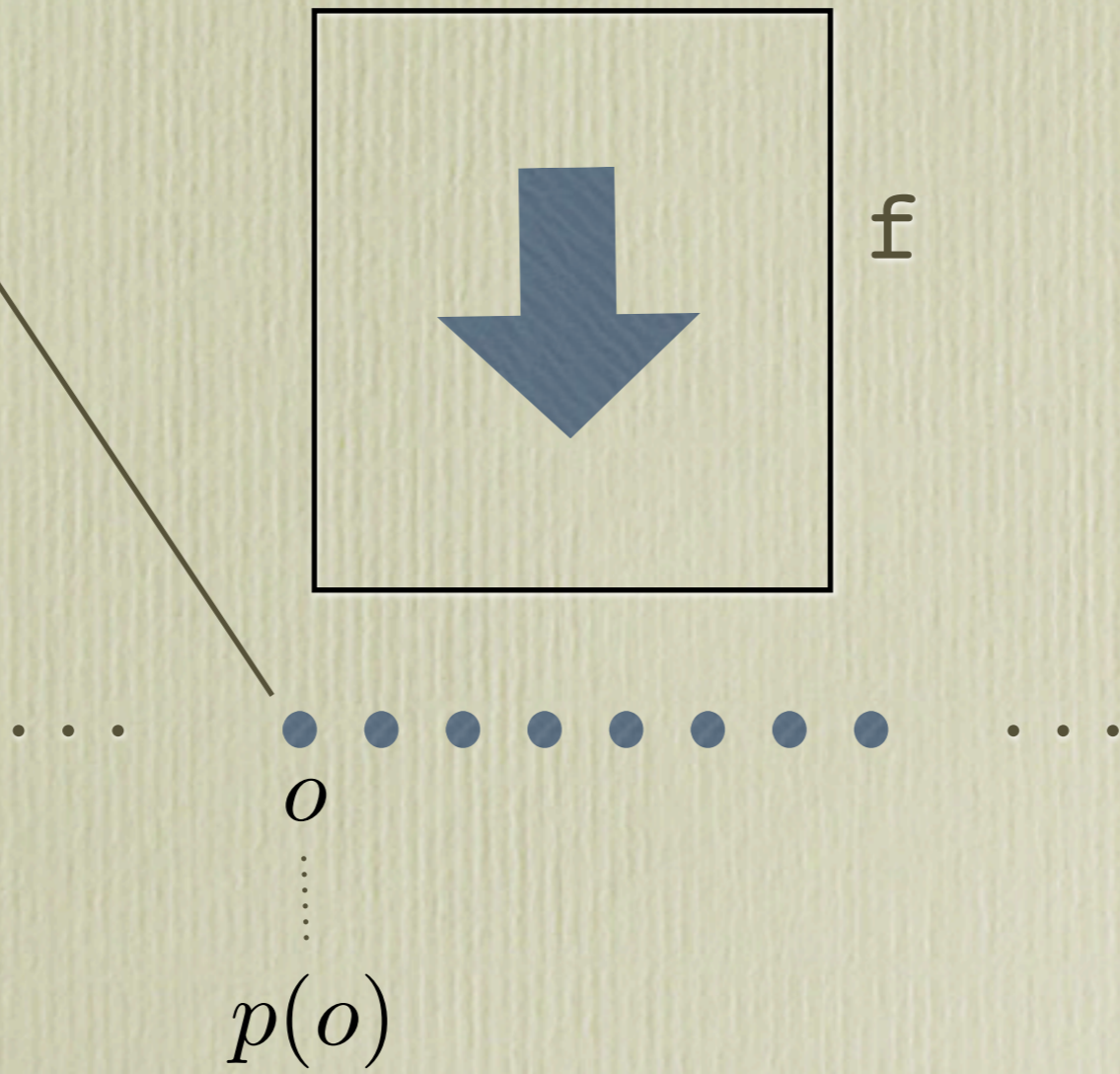
where $[[P]]I = O$

“Simple!”

We call this quantity **Squeeziness**.

$$Sq(f) = \mathcal{H}(I) - \mathcal{H}(O) = \sum_{o \in O} p(o) \mathcal{H}(f^{-1}o)$$

via the partition property



it's not DRR

- Squeeziness is not a refinement of DRR (and vice versa).
- DRR is a cruder measure than Squeeziness and makes fewer distinctions.
- orderings they produce on (f, I) pairs are inconsistent.

the likelihood of collisions

assume uniform distribution on \mathbb{I}

$$|f^{-1}o_i|$$

$$PColl(f) = \sum_{i=1}^n \frac{m_i * (m_i - 1)}{d * (d - 1)}$$

$$|\mathbb{I}|$$

Relationship between Squeeziness and PColl not monotonic

Pearson coefficient

Domain size	Max sub	Corr with Sq	Corr with DRR
1.00E+05	200	0.981	0.849
1.00E+05	200	0.986	0.889
1.00E+05	2000	0.981	0.849
1.00E+05	2000	0.986	0.889
1.00E+06	200	0.971	0.748
1.00E+06	200	0.964	0.686
1.00E+07	200	0.968	0.645
1.00E+07	200	0.975	0.606
1.00E+08	200	0.978	0.584
1.00E+08	200	0.975	0.668

what can we do with Squeeziness?

- (1) Measure how much Software Under Test is inclined to fault masking (not so helpful . . .)
- (2) Improve test set selection?

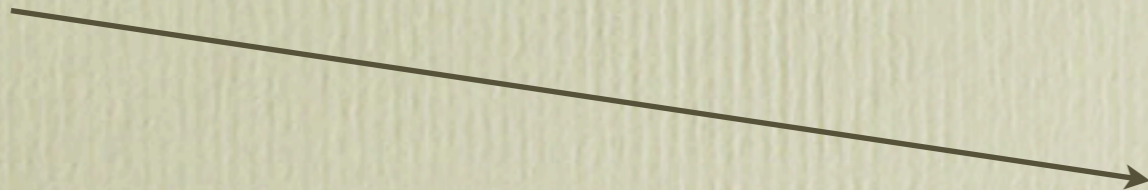
test suite selection

- current “standard” for white box testing is structural coverage: statements, branches, etc.
- limited relationship between coverage and test suite effectiveness, e.g. [Cai and Lyu. A-MOST 2005] plus other papers

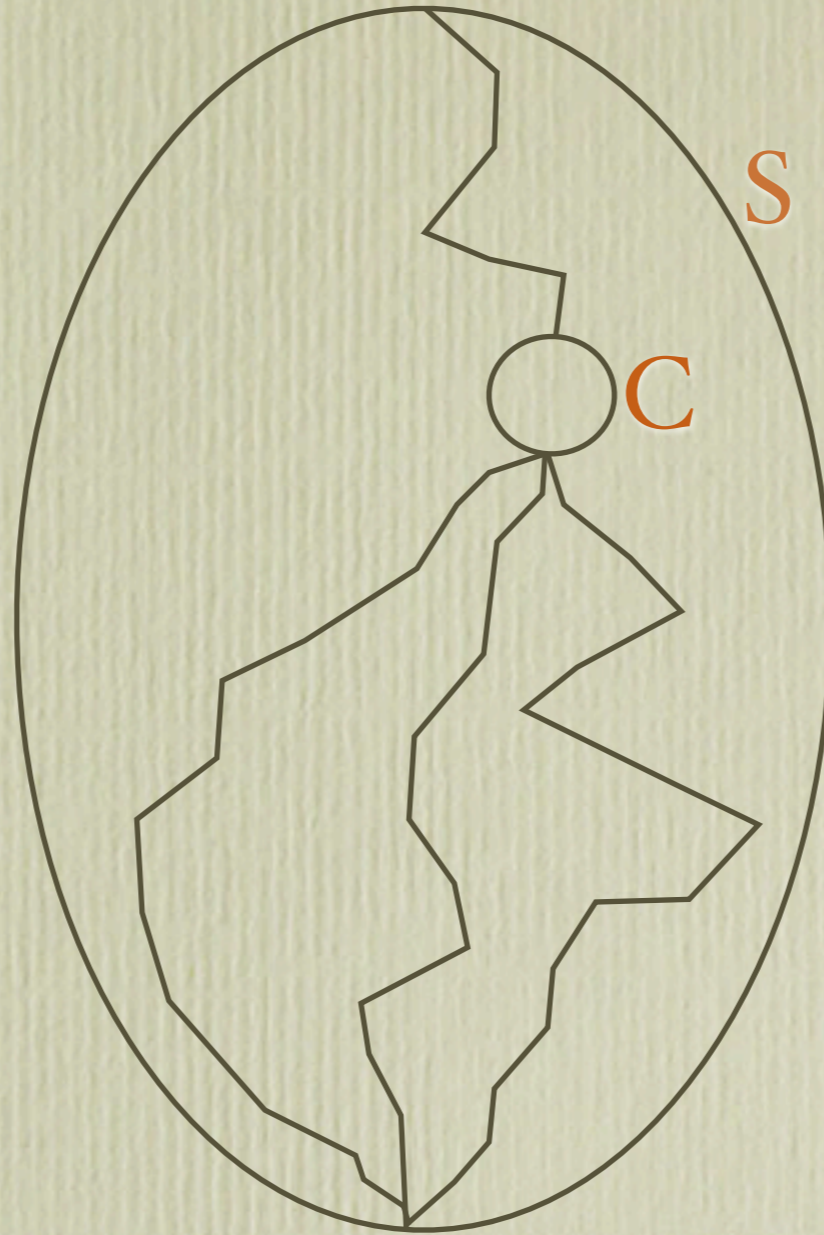
Use covering paths
to generate tests



Pick a less
Squeezy path



Reduce possible
fault masking



probability distributions

How can developers know the random variable in inputs?

(1) Maximum Entropy Principle (= Uniform distribution)

$$Sq(f) = \frac{1}{|I|} \sum_{o \in O} |f^{-1}o| \log_2(|f^{-1}o|)$$

(2) Maximum Squeeziness: $Sq(f) = \log_2|f^{-1}o'|$

(3) Wes Weimar: estimating path execution frequency statically

current research

- experimental validation of post structural element path selection using a mutation testing approach
- theory of probabilistic testing
- program analyses to estimate Squeeziness
- relationship to mutation testing, SBT
- **position paper:** Clark and Hierons. Squeeziness: An Information Theoretic Measure for Avoiding Fault Masking. Accepted for publication in Information Processing Letters

Questions?